Query-Answer Causality in Databases: Abductive Diagnosis and View-Updates

Babak Salimi & Leopoldo Bertossi Carleton University School of Computer Science Ottawa, Canada

• Causality appears at the foundations of many scientific disciplines

- Causality appears at the foundations of many scientific disciplines
- Want to represent and compute causality, to deal with the uncertainty in data and knowledge



- Causality appears at the foundations of many scientific disciplines
- Want to represent and compute causality, to deal with the uncertainty in data and knowledge
- In data management, we need to understand and compute why certain (query) results are obtained or not

- Causality appears at the foundations of many scientific disciplines
- Want to represent and compute causality, to deal with the uncertainty in data and knowledge
- In data management, we need to understand and compute why certain (query) results are obtained or not
- Why certain natural semantic conditions are not satisfied

- Causality appears at the foundations of many scientific disciplines
- Want to represent and compute causality, to deal with the uncertainty in data and knowledge
- In data management, we need to understand and compute why certain (query) results are obtained or not
- Why certain natural semantic conditions are not satisfied
- A DB system should provide explanations

- Causality appears at the foundations of many scientific disciplines
- Want to represent and compute causality, to deal with the uncertainty in data and knowledge
- In data management, we need to understand and compute why certain (query) results are obtained or not
- Why certain natural semantic conditions are not satisfied
- A DB system should provide explanations
- To understand/explore the data or reconsider the query

- Causality appears at the foundations of many scientific disciplines
- Want to represent and compute causality, to deal with the uncertainty in data and knowledge
- In data management, we need to understand and compute why certain (query) results are obtained or not
- Why certain natural semantic conditions are not satisfied
- A DB system should provide explanations
- To understand/explore the data or reconsider the query
- Enhancing the value of extracted data



- Causality appears at the foundations of many scientific disciplines
- Want to represent and compute causality, to deal with the uncertainty in data and knowledge
- In data management, we need to understand and compute why certain (query) results are obtained or not
- Why certain natural semantic conditions are not satisfied
- A DB system should provide explanations
- To understand/explore the data or reconsider the query
- Enhancing the value of extracted data
- Crucial for understanding massive volumes of data!

 Our current research is motivated by trying to understand causality in data management from different perspectives



- Our current research is motivated by trying to understand causality in data management from different perspectives
- We have already established interesting and fruitful connections among three forms of reasoning: (Salimi & Bertossi, ICDT 2015)

- Our current research is motivated by trying to understand causality in data management from different perspectives
- We have already established interesting and fruitful connections among three forms of reasoning: (Salimi & Bertossi, ICDT 2015)
 - inferring causes from databases

- Our current research is motivated by trying to understand causality in data management from different perspectives
- We have already established interesting and fruitful connections among three forms of reasoning: (Salimi & Bertossi, ICDT 2015)
 - inferring causes from databases
 - database repairs and consistent query answering (CQA)



- Our current research is motivated by trying to understand causality in data management from different perspectives
- We have already established interesting and fruitful connections among three forms of reasoning: (Salimi & Bertossi, ICDT 2015)
 - inferring causes from databases
 - database repairs and consistent query answering (CQA)
 - consistency-based diagnoses

- Our current research is motivated by trying to understand causality in data management from different perspectives
- We have already established interesting and fruitful connections among three forms of reasoning: (Salimi & Bertossi, ICDT 2015)
 - inferring causes from databases
 - database repairs and consistent query answering (CQA)
 - consistency-based diagnoses
- In this work we add:

- Our current research is motivated by trying to understand causality in data management from different perspectives
- We have already established interesting and fruitful connections among three forms of reasoning: (Salimi & Bertossi, ICDT 2015)
 - inferring causes from databases
 - database repairs and consistent query answering (CQA)
 - consistency-based diagnoses
- In this work we add:
 - view updates (database updates through views)

- Our current research is motivated by trying to understand causality in data management from different perspectives
- We have already established interesting and fruitful connections among three forms of reasoning: (Salimi & Bertossi, ICDT 2015)
 - inferring causes from databases
 - database repairs and consistent query answering (CQA)
 - consistency-based diagnoses
- In this work we add:
 - view updates (database updates through views)
 - abductive diagnosis



• Those reasoning problems share some commonalities:



- Those reasoning problems share some commonalities:
 - They reflect some sort of uncertain information



- Those reasoning problems share some commonalities:
 - They reflect some sort of uncertain information
 - They are *non-monotonic*

- Those reasoning problems share some commonalities:
 - They reflect some sort of uncertain information
 - They are *non-monotonic*
 - They are forms of reverse data transformation
 To reason "diagnostically" about an output (observation)



- Those reasoning problems share some commonalities:
 - They reflect some sort of *uncertain information*
 - They are non-monotonic
 - They are forms of reverse data transformation
 To reason "diagnostically" about an output (observation)



 These problems have been classified under reverse data management (Meliou et al., VLDB 2011)

```
• Assume D = D^n \cup D^x

D^n endogenous tuples (candidate causes)

D^x exogenous tuples
```

- Assume D = Dⁿ ∪ D^x
 Dⁿ endogenous tuples (candidate causes)
 D^x exogenous tuples
- $\tau \in D^n$ is counterfactual cause for answer \bar{a} to monotone query $\mathcal{Q}(\bar{x})$ in D, if $D \models \mathcal{Q}(\bar{a})$ and $D \setminus \{\tau\} \not\models \mathcal{Q}(\bar{a})$

- Assume D = Dⁿ ∪ D^x
 Dⁿ endogenous tuples (candidate causes)
 D^x exogenous tuples
- $\tau \in D^n$ is counterfactual cause for answer \bar{a} to monotone query $\mathcal{Q}(\bar{x})$ in D, if $D \models \mathcal{Q}(\bar{a})$ and $D \setminus \{\tau\} \not\models \mathcal{Q}(\bar{a})$
- $\tau \in D^n$ is an actual cause for \bar{a} if there is $\Gamma \subseteq D^n$, a contingency set, such that τ is a counterfactual cause for \bar{a} in $D \setminus \Gamma$

A notion of causality-based explanation for a query result was introduced in (Meliou et al., VLDB 2010)

- Assume D = Dⁿ ∪ D^x
 Dⁿ endogenous tuples (candidate causes)
 D^x exogenous tuples
- $\tau \in D^n$ is counterfactual cause for answer \bar{a} to monotone query $\mathcal{Q}(\bar{x})$ in D, if $D \models \mathcal{Q}(\bar{a})$ and $D \setminus \{\tau\} \not\models \mathcal{Q}(\bar{a})$
- $\tau \in D^n$ is an actual cause for \bar{a} if there is $\Gamma \subseteq D^n$, a contingency set, such that τ is a counterfactual cause for \bar{a} in $D \setminus \Gamma$

Based on Halpern & Pearl's actual causation w/deletion as "intervention" (Halpern & Pearl, 2001, 2005)

Causal Responsibility

Causal responsibility reflects *relative degree of causality* of a tuple for a query result (Meliou et al., VLDB 2010)

• The responsibility of an actual cause τ for \bar{a} :

$$ho_{\!{\scriptscriptstyle D},{\scriptscriptstyle {\mathcal Q}(ar{ar{m{z}}})}}\!\!\left(au
ight)=rac{1}{|\Gamma|+1}$$

 $|\Gamma|=$ size of smallest contingency set for au

Causal Responsibility

Causal responsibility reflects *relative degree of causality* of a tuple for a query result (Meliou et al., VLDB 2010)

• The responsibility of an actual cause τ for \bar{a} :

$$ho_{\!{\scriptscriptstyle D},{\scriptscriptstyle {\mathcal Q}(ar{ar{m{z}}})}}\!\!\left(au
ight)=rac{1}{|\Gamma|+1}$$

 $|\Gamma|=$ size of smallest contingency set for au

Tuples with higher responsibility tend to provide more interesting explanations for query results

Based on (Chockler and Halpern, 2004)

Example: Database D

Author	AuName	Journal
	Joe	TKDE
	John	TKDE
	Tom	TKDE
	John	TODS

Journal	Journal	Topic	#Paper
	TKDE	XML	30
	TKDE	CUBE	30
	TODS	XML	30

Example: Database D

Author	AuName	Journal
	Joe	TKDE
	John	TKDE
	Tom	TKDE
	John	TODS

Journal	Journal	Topic	#Paper
	TKDE	XML	30
	TKDE	CUBE	30
	TODS	XML	30

Conjunctive query Q(x, y):

```
Ans(AuName, Topic) \leftarrow Author(AuName, Journal),

Journal(Journal, Topic, \#Paper),
```

Example: Database D

Author	AuName	Journal
	Joe	TKDE
	John	TKDE
	Tom	TKDE
	John	TODS

Journal	Journal	Topic	#Paper
	TKDE	XML	30
	TKDE	CUBE	30
	TODS	XML	30

Conjunctive query Q(x, y):

Q(D)	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE

Example: (cont.)

(John, XML): unexpected answer

What are the causes ?

Author	AuName	Journal	Journal	Journal	Topic	#Paper
	Joe John Tom John	TKDE TKDE TKDE TODS	Journal	TKDE TKDE TODS	XML CUBE XML	30 30 30

 $Ans_{\mathcal{Q}}(AuName, Topic) \leftarrow Author(AuName, Journal),$ Journal(Journal, Topic, #Paper),

$\mathcal{Q}(D)$	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE

Example: (cont.)

(John, XML): unexpected answer

What are the causes?

 $\tau = Author(John, TODS)$ is an actual cause, with contingency sets

 $\Gamma_1 = \{ Author(John, TKDE) \}$ and

Journal(Journal, Topic, #Paper), $\Gamma_2 = \{Journal(TKDE, XML, 30)\}$

	Joe John Tom John	TKDE TKDE TKDE TODS	Journal	TKDE TKDE TODS	XML CUBE XML	30 30 30 30	
Anc - (AuNama	Tonic)	,	Author	AuNama	lour	I'

Q(D)	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE

Example: (cont.)

(John, XML): unexpected answer

What are the causes?

 $\tau = Author(John, TODS)$ is an actual cause, with contingency sets

 $\Gamma_1 {=} \{ \text{ } \textit{Author(John, TKDE)} \ \} \ \text{ and } \\$

 \leftarrow Author(AuName, Journal), Journal(Journal, Topic, #Paper), $\Gamma_2 = \{Journal(TKDE, XML, 30)\}$

 $\rho(\tau) = \frac{1}{2}$

Author	AuName	Journal			Topic	#Paper
	Joe	TKDE	Journal	Journal TKDE	XML.	#Paper 30
	John	TKDE		TKDE	CUBE	30
	Tom John	TKDE		TODS	XML	30

 $Ans_{\mathcal{Q}}(AuName, Topic) \leftarrow Author(AuName, Journal),$ Journal(Journal, Topic, #Paper)

$\mathcal{Q}(D)$	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE

Example: (cont.)

(John, XML): unexpected answer

What are the causes?

 $\tau = Author(John, TODS)$ is an actual cause, with contingency sets

 $\Gamma_1 = \{ Author(John, TKDE) \}$ and

 \leftarrow Author(AuName, Journal),
Journal(Journal, Topic, #Paper), $\Gamma_2 = \{Journal(TKDE, XML, 30)\}$

John TKDE Tom TKDE John TODS		TKDE	CUBE	30 30	
Anso(AuName, Topic)	←	Author(AuNam	e. Jour	nal

Answers:

$\mathcal{Q}(D)$	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE
	John	CORF

$$\rho(\tau) = \frac{1}{2}$$

Journal(TKDE, XML, 30), Author(John, TKDE), Journal(TODS,XML, 30) are also actual causes with responsibility $\frac{1}{2}$

 Intuitively, an abductive explanation of an observation is a formula that, together with the background logical theory, entails the observation

- Intuitively, an abductive explanation of an observation is a formula that, together with the background logical theory, entails the observation
- The background theory specifies the system being diagnosed

- Intuitively, an abductive explanation of an observation is a formula that, together with the background logical theory, entails the observation
- The background theory specifies the system being diagnosed
- It is possible to specify systems with *Datalog programs*



- Intuitively, an abductive explanation of an observation is a formula that, together with the background logical theory, entails the observation
- The background theory specifies the system being diagnosed
- It is possible to specify systems with Datalog programs
 - An example of (recursive) Datalog program:

```
parent(mary, john)

parent(mary, john)

ancestor(X, Y) \leftarrow parent(X, Y)

ancestor(X, Y) \leftarrow parent(X, Z), ancestor(Z, Y)
```

- Intuitively, an abductive explanation of an observation is a formula that, together with the background logical theory, entails the observation
- The background theory specifies the system being diagnosed
- It is possible to specify systems with Datalog programs
 - An example of (recursive) Datalog program:

```
parent(mary, john)

parent(mary, john)

ancestor(X, Y) \leftarrow parent(X, Y)

ancestor(X, Y) \leftarrow parent(X, Z), ancestor(Z, Y)
```

 Datalog programs define monotone queries and the notion of cause can be applied as above

• A Datalog abduction problem (DAP) is of the form $\mathcal{AP} = \langle \Pi, E, Hyp, Obs \rangle$,

• A *Datalog abduction problem* (DAP) is of the form

$$\mathcal{AP} = \langle \Pi, E, Hyp, Obs \rangle$$
,

(a) Π: set of Datalog rules

```
\mathcal{AP} = \langle \Pi, E, Hyp, Obs \rangle,
```

- (a) Π: set of Datalog rules
- (b) E: finite set of ground atoms (extensional database)

```
\mathcal{AP} = \langle \Pi, E, Hyp, Obs \rangle,
```

- (a) Π: set of Datalog rules
- (b) E: finite set of ground atoms (extensional database)
- (c) Hyp: finite set of ground atoms, the abducible atoms

```
\mathcal{AP} = \langle \Pi, E, Hyp, Obs \rangle,
```

- (a) Π: set of Datalog rules
- (b) E: finite set of ground atoms (extensional database)
- (c) Hyp: finite set of ground atoms, the abducible atoms
- (d) *Obs*: observation, a finite conjunction of ground atoms with $\Pi \cup E \cup Hyp \models Obs$

```
\mathcal{AP} = \langle \Pi, E, Hyp, Obs \rangle,
```

- (a) Π: set of Datalog rules
- (b) E: finite set of ground atoms (extensional database)
- (c) Hyp: finite set of ground atoms, the abducible atoms
- (d) *Obs*: observation, a finite conjunction of ground atoms with $\Pi \cup E \cup Hyp \models Obs$
- The abduction problem is about computing a subset-minimal $\Delta \subseteq Hyp$, such that $\Pi \cup E \cup \Delta \models Obs$

```
\mathcal{AP} = \langle \Pi, E, Hyp, Obs \rangle,
```

- (a) Π: set of Datalog rules
- (b) E: finite set of ground atoms (extensional database)
- (c) Hyp: finite set of ground atoms, the abducible atoms
- (d) *Obs*: observation, a finite conjunction of ground atoms with $\Pi \cup E \cup Hyp \models Obs$
- The abduction problem is about computing a subset-minimal $\Delta \subseteq Hyp$, such that $\Pi \cup E \cup \Delta \models Obs$
- Relevance Problem: Deciding if $h \in Hyp$ belongs to some abductive diagnosis NP-complete! (in $|\mathcal{AP}|$)

• View: virtual table defined by a query, e.g. a conjunctive query

- View: virtual table defined by a query, e.g. a conjunctive query
- *View-update problem:* update the DB as propagation of changes on views, i.e. from views to base relations

- View: virtual table defined by a query, e.g. a conjunctive query
- View-update problem: update the DB as propagation of changes on views, i.e. from views to base relations
- When view updates are tuple deletions, this is delete-propagation problem

- View: virtual table defined by a query, e.g. a conjunctive query
- *View-update problem:* update the DB as propagation of changes on views, i.e. from views to base relations
- When view updates are tuple deletions, this is delete-propagation problem
 - How to delete tuples from the database, so that an undesired tuple disappears from the view
 Several variants:

- View: virtual table defined by a query, e.g. a conjunctive query
- *View-update problem:* update the DB as propagation of changes on views, i.e. from views to base relations
- When view updates are tuple deletions, this is <u>delete-propagation</u> problem
 - How to delete tuples from the database, so that an undesired tuple disappears from the view
 Several variants:
 - delete subset-minimal set of tuples from source (minimal source side-effect problem)

- View: virtual table defined by a query, e.g. a conjunctive query
- View-update problem: update the DB as propagation of changes on views, i.e. from views to base relations
- When view updates are tuple deletions, this is delete-propagation problem
 - How to delete tuples from the database, so that an undesired tuple disappears from the view
 Several variants:
 - delete subset-minimal set of tuples from source (minimal source side-effect problem)
 - delete minimum number of tuples from source (minimum source side-effect problem)

Author	AuName Joe John	Journal TKDE TKDE	Journal	Journal TKDE TKDE	Topic XML CUBE	#Paper 30 30	
	Tom John	TODS		TODS	XML	30	
Ans _Q (.	AuName	, Topic)				ne, Journ pic, #Pa	- "

For view V defined by above query:

V(D)	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE

Don't want (John, XML) in the view extension

Author	Joe John Tom John	Journal TKDE TKDE TKDE TCDS	Journal	Journal TKDE TKDE TODS	Topic XML CUBE XML	#Paper 30 30 30 30	
Ans _Q (AuName	, Topic)				ne, Journ pic, #Pap	- "

For view V defined by above query:

V(D)	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE

Don't want (John, XML) in the view extension

Minimal source side-effect solutions:

```
\begin{split} & \rho_1 = \{ \ \textit{Author}(\mathsf{John}, \, \mathsf{TODS}), \ \textit{Journal}(\mathsf{TODS}, \, \mathsf{XML}, 30) \} \\ & \rho_2 = \{ \ \textit{Author}(\mathsf{John}, \, \mathsf{TODS}), \ \textit{Author}(\mathsf{John}, \, \mathsf{TKDE}) \} \\ & \rho_3 = \{ \ \textit{Author}(\mathsf{John}, \, \mathsf{TDK}), \ \textit{Journal}(\mathsf{TODS}, \, \mathsf{XML}, 30) \} \\ & \rho_4 = \{ \ \textit{Journal}(\mathsf{TODS}, \, \mathsf{XML}, 30), \ \textit{Journal}(\mathsf{John}, \, \mathsf{TKDE}, 30) \} \end{split}
```

Ans _Q (AuName	, Topic)				ne, Journa pic, #Pa _l	
Author	Joe John Tom John	TKDE TKDE TKDE TKDE TODS	Journal	Journal TKDE TKDE TODS	Topic XML CUBE XML	#Paper 30 30 30 30	

For view V defined by above query:

V(D)	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE

Don't want (John, XML) in the view extension

Minimal source side-effect solutions:

```
p_1=\{ Author(John, TODS), Journal(TODS, XML,30)}

p_2=\{ Author(John, TODS), Author(John, TKDE)}

p_3=\{ Author(John, TDK), Journal(TODS, XML,30)}

p_4=\{ Journal(TODS, XML,30), Journal(John, TKDE,30)}
```

Also solutions to minimum source sideeffect problem

Ans _Q (AuName	, Topic)				ne, Journa pic, #Pa _l	
Author	Joe John Tom John	TKDE TKDE TKDE TKDE TODS	Journal	Journal TKDE TKDE TODS	Topic XML CUBE XML	#Paper 30 30 30 30	

For view V defined by above query:

V(D)	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE

Don't want (John, XML) in the view extension

Minimal source side-effect solutions:

```
p_1=\{ Author(John, TODS), Journal(TODS, XML,30)}

p_2=\{ Author(John, TODS), Author(John, TKDE)}

p_3=\{ Author(John, TDK), Journal(TODS, XML,30)}

p_4=\{ Journal(TODS, XML,30), Journal(John, TKDE,30)}
```

Also solutions to minimum source sideeffect problem

Author	AuName	Journal	Journal	Journal	Topic	#Paper	
	Joe John Tom John	TKDE TKDE TKDE TODS		TKDE TKDE TODS	XML CUBE XML	30 30 30	
$Ans_{\mathcal{Q}}($	AuName	, Topic)				ne, Journ	- "
			Journ	ai (Jouri	iai, roj	JIC, #Fa	per)

For view V defined by above query:

V(D)	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE

Don't want (John, XML) in the view extension

Minimal source side-effect solutions:

```
p_1=\{ Author(John, TODS), Journal(TODS, XML,30)}

p_2=\{ Author(John, TODS), Author(John, TKDE)}

p_3=\{ Author(John, TDK), Journal(TODS, XML,30)}

p_4=\{ Journal(TODS, XML,30), Journal(John, TKDE,30)}
```

Also solutions to minimum source sideeffect problem

Notice similarities with causality notions!

• In this work we established precise connections between:

- In this work we established precise connections between:
 - Causality: a recent problem in databases

- In this work we established precise connections between:
 - Causality: a recent problem in databases
 - Datalog abduction
 Abductive diagnosis is an important problem in KR
 Also of interest in databases

- In this work we established precise connections between:
 - Causality: a recent problem in databases
 - Datalog abduction
 Abductive diagnosis is an important problem in KR
 Also of interest in databases
 - Deletion-based view updates (interesting case for monotone views)
 View updates is a classical and important problem in databases

- In this work we established precise connections between:
 - Causality: a recent problem in databases
 - Datalog abduction
 Abductive diagnosis is an important problem in KR
 Also of interest in databases
 - Deletion-based view updates (interesting case for monotone views)
 View updates is a classical and important problem in databases
- Characterizations of each in terms of the others (abduction vs. view-updates has been investigated before)

- In this work we established precise connections between:
 - Causality: a recent problem in databases
 - Datalog abduction
 Abductive diagnosis is an important problem in KR
 Also of interest in databases
 - Deletion-based view updates (interesting case for monotone views)
 View updates is a classical and important problem in databases
- Characterizations of each in terms of the others (abduction vs. view-updates has been investigated before)
- Known complexity results for some of them are applied to the others, by reduction

 Query-answer causality in databases has a close connection with Datalog abduction

Causal and abductive entailments coincide

 Query-answer causality in databases has a close connection with Datalog abduction

Causal and abductive entailments coincide

 Actual causes and their responsibility for a Datalog query can obtained from abductive diagnoses of a corresponding Datalog abduction problem

 Query-answer causality in databases has a close connection with Datalog abduction

Causal and abductive entailments coincide

- Actual causes and their responsibility for a Datalog query can obtained from abductive diagnoses of a corresponding Datalog abduction problem
- Relevant hypothesis for a Datalog abduction problem can be obtained from actual causes for a corresponding causality problem

Example: Instance D (no exogenous tuples) and Boolean query

$$\Pi$$
: ans $\leftarrow R(x, y), S(y)$,

R	Х	Υ
	a_1	a 4
	a 2	a_1
	a 3	<i>a</i> ₃

S	Χ
	a_1
	a ₂
	<i>a</i> ₃

Example: Instance D (no exogenous tuples) and Boolean query

$$\Pi$$
: ans $\leftarrow R(x,y), S(y),$

R	Χ	Υ
	a_1	<i>a</i> ₄
	a ₂	<i>a</i> ₁
	<i>a</i> ₃	<i>a</i> ₃

Tuples $S(a_1)$, $R(a_2, a_1)$, $S(a_3)$ and $R(a_3, a_3)$ are actual causes for ans, with responsibility $\frac{1}{2}$

Example: (as above) D (no exogenous tuple) Π : $ans \leftarrow R(x, y), S(y)$

R	Χ	Υ
	<i>a</i> ₁	a 4
	<i>a</i> ₂	a_1
	<i>a</i> ₃	<i>a</i> ₃

Example: (as above) D (no exogenous tuple) Π : $ans \leftarrow R(x, y), S(y)$

R	Χ	Υ
	<i>a</i> ₁	a 4
	a ₂	a_1
	<i>a</i> ₃	<i>a</i> ₃

Consider the Datalog abduction problem: $\mathcal{AP}^c = \langle \Pi, \emptyset, D, \{ans\} \rangle$

Example: (as above) D (no exogenous tuple) Π : $ans \leftarrow R(x, y), S(y)$

R	Х	Υ
	a_1	a 4
	a ₂	a_1
	<i>a</i> ₃	<i>a</i> ₃

Consider the Datalog abduction problem: $\mathcal{AP}^c = \langle \Pi, \emptyset, D, \{ans\} \rangle$

 \mathcal{AP}^{c} has two (subset-minimal) abductive diagnosis:

Example: (as above) D (no exogenous tuple) Π : $ans \leftarrow R(x, y), S(y)$

R	Χ	Υ
	<i>a</i> ₁	a 4
	a ₂	a_1
	<i>a</i> ₃	<i>a</i> ₃

Consider the Datalog abduction problem: $\mathcal{AP}^c = \langle \Pi, \emptyset, D, \{ans\} \rangle$

 \mathcal{AP}^{c} has two (subset-minimal) abductive diagnosis:

•
$$\Delta_1 = \{S(a_1), R(a_2, a_1)\}$$

Example: (as above) D (no exogenous tuple) Π : $ans \leftarrow R(x, y), S(y)$

R	Χ	Υ
	a_1	a 4
	a ₂	a_1
	<i>a</i> ₃	<i>a</i> ₃

Consider the Datalog abduction problem: $\mathcal{AP}^c = \langle \Pi, \emptyset, D, \{ans\} \rangle$

 \mathcal{AP}^{c} has two (subset-minimal) abductive diagnosis:

- $\Delta_1 = \{ S(a_1), R(a_2, a_1) \}$
- $\bullet \ \Delta_2 = \{S(a_3), R(a_3, a_3)\}\$

Example: (as above) D (no exogenous tuple) Π : $ans \leftarrow R(x, y), S(y)$

R	Χ	Υ
	a_1	a 4
	a ₂	a_1
	a 3	a 3

Consider the Datalog abduction problem: $\mathcal{AP}^c = \langle \Pi, \emptyset, D, \{ans\} \rangle$

 \mathcal{AP}^{c} has two (subset-minimal) abductive diagnosis:

- $\Delta_1 = \{ S(a_1), R(a_2, a_1) \}$

 $S(a_3)$, $R(a_3, a_3)$, $S(a_1)$ and $R(a_2, a_1)$ are relevant hypothesis

Example: (as above) D (no exogenous tuple) Π : $ans \leftarrow R(x, y), S(y)$

R	Χ	Υ
	a_1	a 4
	a ₂	a_1
	<i>a</i> ₃	<i>a</i> ₃

Consider the Datalog abduction problem: $\mathcal{AP}^c = \langle \Pi, \emptyset, D, \{ans\} \rangle$

 \mathcal{AP}^{c} has two (subset-minimal) abductive diagnosis:

- $\bullet \ \Delta_1 = \{S(a_1), R(a_2, a_1)\}$
- $\bullet \ \Delta_2 = \{S(a_3), R(a_3, a_3)\}$

 $S(a_3)$, $R(a_3, a_3)$, $S(a_1)$ and $R(a_2, a_1)$ are relevant hypothesis

And also the actual causes for ans!



Example: (as above)

R	Х	Υ
	a_1	a 4
	a ₂	a_1
	<i>a</i> ₃	<i>a</i> ₃

$$\mathcal{AP}^c$$
 as before, $\Delta_1 = \{S(a_1), R(a_2, a_1)\}, \ \Delta_2 = \{S(a_3), R(a_3, a_3)\}$

$$\frac{\mathsf{Example}:}{\mathsf{R}}$$
 (as above)

R	Χ	Υ
	a_1	a 4
	a ₂	a_1
	<i>a</i> ₃	<i>a</i> ₃

$$\mathcal{AP}^c$$
 as before, $\Delta_1 = \{S(a_1), R(a_2, a_1)\}, \ \Delta_2 = \{S(a_3), R(a_3, a_3)\}$

To obtain responsibilities, we compute *necessary-hypothesis sets* of \mathcal{AP}^c

R	Χ	Υ
	<i>a</i> ₁	a 4
	a ₂	a_1
	<i>a</i> ₃	<i>a</i> ₃

$$\mathcal{AP}^c$$
 as before, $\Delta_1 = \{S(a_1), R(a_2, a_1)\}, \ \Delta_2 = \{S(a_3), R(a_3, a_3)\}$

To obtain responsibilities, we compute *necessary-hypothesis sets* of \mathcal{AP}^c

W/O them in Hyp there is no abductive diagnosis

Example: (as above) _

R	Χ	Υ
	<i>a</i> ₁	a 4
	a ₂	a_1
	<i>a</i> ₃	<i>a</i> ₃

$$\mathcal{AP}^c$$
 as before, $\Delta_1 = \{S(a_1), R(a_2, a_1)\}, \ \Delta_2 = \{S(a_3), R(a_3, a_3)\}$

To obtain responsibilities, we compute *necessary-hypothesis sets* of \mathcal{AP}^c

W/O them in Hyp there is no abductive diagnosis

- $N_1 = \{S(a_1), S(a_3)\}, \quad N_2 = \{S(a_3), R(a_2, a_1)\}$
- $N_3 = \{S(a_1), R(a_3, a_3)\}, N_4 = \{R(a_2, a_1), R(a_3, a_3)\}$

R	Χ	Υ
	<i>a</i> ₁	<i>a</i> ₄
	a ₂	a_1
	<i>a</i> ₃	<i>a</i> ₃

$$\mathcal{AP}^c$$
 as before, $\Delta_1 = \{S(a_1), R(a_2, a_1)\}, \ \Delta_2 = \{S(a_3), R(a_3, a_3)\}$

To obtain responsibilities, we compute *necessary-hypothesis sets* of \mathcal{AP}^c

W/O them in Hyp there is no abductive diagnosis

- $N_1 = \{S(a_1), S(a_3)\}, N_2 = \{S(a_3), R(a_2, a_1)\}$
- $N_3 = \{S(a_1), R(a_3, a_3)\}, N_4 = \{R(a_2, a_1), R(a_3, a_3)\}$

It holds: The responsibilities of causal tuples τ for ans are $\frac{1}{|N|}$, with N a minimum-cardinality necessary-hypothesis set, and $\tau \in N$

Example: (as above)
$$\begin{array}{c|c}
\hline
R & X & Y \\
\hline
a_1 & a_4 \\
a_2 & a_1 \\
a_3 & a_3
\end{array}$$

$$\mathcal{AP}^c$$
 as before, $\Delta_1 = \{S(a_1), R(a_2, a_1)\}, \ \Delta_2 = \{S(a_3), R(a_3, a_3)\}$

To obtain responsibilities, we compute necessary-hypothesis sets of \mathcal{AP}^c

W/O them in Hyp there is no abductive diagnosis

- $N_1 = \{S(a_1), S(a_3)\}, N_2 = \{S(a_3), R(a_2, a_1)\}$
- $N_3 = \{S(a_1), R(a_3, a_3)\}, N_4 = \{R(a_2, a_1), R(a_3, a_3)\}$

It holds: The responsibilities of causal tuples τ for ans are $\frac{1}{|N|}$, with N a minimum-cardinality necessary-hypothesis set, and $\tau \in N$

Again, all causes have responsibility $\frac{1}{2}$

• These connections enable mutual applications



- These connections enable mutual applications
- For example, for Datalog queries, deciding if a tuple is a cause is NP-complete (combined complexity)

- These connections enable mutual applications
- For example, for Datalog queries, deciding if a tuple is a cause is *NP*-complete (combined complexity)
 - Obtained from complexity of relevance problem for Datalog abduction

- These connections enable mutual applications
- For example, for Datalog queries, deciding if a tuple is a cause is *NP*-complete (combined complexity)
 - Obtained from complexity of relevance problem for Datalog abduction
 - Notice that for unions of Boolean conjunctive queries the problem is tractable in data (Salimi & Bertossi, ICDT 2015)

- These connections enable mutual applications
- For example, for Datalog queries, deciding if a tuple is a cause is *NP*-complete (combined complexity)
 - Obtained from complexity of relevance problem for Datalog abduction
 - Notice that for unions of Boolean conjunctive queries the problem is tractable in data (Salimi & Bertossi, ICDT 2015)
- We identify tractable classes of deciding causality for Datalog queries and instances

- These connections enable mutual applications
- For example, for Datalog queries, deciding if a tuple is a cause is *NP*-complete (combined complexity)
 - Obtained from complexity of relevance problem for Datalog abduction
 - Notice that for unions of Boolean conjunctive queries the problem is tractable in data (Salimi & Bertossi, ICDT 2015)
- We identify tractable classes of deciding causality for Datalog queries and instances
 - Obtained from tractable cases of Datalog abduction
 Guarded programs, bounded tree-width instances
 (Gottlob, Pichler & Wei, 2010)

 As suggested above, there is a close relationship between query causality and delete-propagation for view-updates

- As suggested above, there is a close relationship between query causality and delete-propagation for view-updates
- Want to propagate deletion a tuple \bar{v} from view V to underlying D?

- As suggested above, there is a close relationship between query causality and delete-propagation for view-updates
- Want to propagate deletion a tuple \bar{v} from view V to underlying D?
- If V defined by monotone query $\mathcal{Q}(\bar{x})$, delete from D the actual causes for answer \bar{v}

- As suggested above, there is a close relationship between query causality and delete-propagation for view-updates
- Want to propagate deletion a tuple \bar{v} from view V to underlying D?
- If V defined by monotone query $\mathcal{Q}(\bar{x})$, delete from D the actual causes for answer \bar{v}
- Actual causes with subset-minimal contingency sets give solutions to minimal source side-effect problem (and viceversa)

- As suggested above, there is a close relationship between query causality and delete-propagation for view-updates
- Want to propagate deletion a tuple \bar{v} from view V to underlying D?
- If V defined by monotone query $\mathcal{Q}(\bar{x})$, delete from D the actual causes for answer \bar{v}
- Actual causes with subset-minimal contingency sets give solutions to minimal source side-effect problem (and viceversa)
- Most responsible actual causes with minimum-cardinality contingency sets give solutions to minimum source side-effect problem (and viceversa)

Example: Again

Author	AuName	Journal	Journal	Journal	Topic	#Paper
	Joe John Tom John	TKDE TKDE TKDE TODS	Journal	TKDE TKDE TODS	XML CUBE XML	30 30 30 30

 $Ans_{\mathcal{Q}}(AuName, Topic) \leftarrow Author(AuName, Journal),$ Journal(Journal, Topic, #Paper),

View:

V(D)	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE

(John, XML) not wanted in the view

Example: Again

Author	Joe TKDI John TKDI Tom TKDI	Journal	Journal	Journal	Topic	#Paper
		TKDE TKDE TODS		TKDE TKDE TODS	XML CUBE XML	30 30 30

 $Ans_{\mathcal{Q}}(AuName, Topic) \leftarrow Author(AuName, Journal),$ Journal(Journal, Topic, #Paper).

Combination for (John, XML) of actual causes/contingency sets

(each element is actual cause and the complement a contingency set)

View:

V(D)	AuName	Topic
	Joe	XML
	Joe	CUBE
	Tom	XML
	Tom	CUBE
	John	XML
	John	CUBE

$$\begin{split} p_1 &= \{ \text{ Author}(\text{John, TODS}), \text{ Journal}(\text{TODS}, \text{XML}, 30) \} \\ p_2 &= \{ \text{ Author}(\text{John, TODS}), \text{ Author}(\text{John, TKDE}) \} \\ p_3 &= \{ \text{ Author}(\text{John, TDK}), \text{ Journal}(\text{TODS}, \text{XML}, 30) \} \\ p_4 &= \{ \text{ Journal}(\text{TODS}, \text{XML}, 30), \text{ Journal}(\text{John, TKDE}, 30) \} \end{split}$$

(John, XML) not wanted in the view

We can take advantage of established connections:

We can take advantage of established connections:

- Computing the size of a solution to a minimum source side-effect problem is FP^{NP(log(n))}-hard in data
 - From complexity of computing most responsible causes in (Salimi and Bertossi ICDT, 2015)
- We identify class of queries for which the minimum source side-effect problem is tractable
 - From the dichotomy result for complexity of responsibility in (Meliou et al., VLDB 2010)

 Causality has been a research subject in AI, Statistics, etc. etc. for many years

- Causality has been a research subject in AI, Statistics, etc. etc. for many years
- Causality in data management (DM) is much newer subject

- Causality has been a research subject in AI, Statistics, etc. etc. for many years
- Causality in data management (DM) is much newer subject
- We have started to scratch the surface (as a DM community)

- Causality has been a research subject in AI, Statistics, etc. etc. for many years
- Causality in data management (DM) is much newer subject
- We have started to scratch the surface (as a DM community)
- Many extensions are possible and necessary, to understand data phenomena in causal terms

- Causality has been a research subject in AI, Statistics, etc. etc. for many years
- Causality in data management (DM) is much newer subject
- We have started to scratch the surface (as a DM community)
- Many extensions are possible and necessary, to understand data phenomena in causal terms
- Interesting results on causality in DM have been obtained

- Causality has been a research subject in AI, Statistics, etc. etc. for many years
- Causality in data management (DM) is much newer subject
- We have started to scratch the surface (as a DM community)
- Many extensions are possible and necessary, to understand data phenomena in causal terms
- Interesting results on causality in DM have been obtained
- Causality in DM is related to many other DM reasoning tasks

- Causality has been a research subject in AI, Statistics, etc. etc. for many years
- Causality in data management (DM) is much newer subject
- We have started to scratch the surface (as a DM community)
- Many extensions are possible and necessary, to understand data phenomena in causal terms
- Interesting results on causality in DM have been obtained
- Causality in DM is related to many other DM reasoning tasks
- Maybe it is -unsurprisingly- an important underlying principle

- Causality has been a research subject in AI, Statistics, etc. etc. for many years
- Causality in data management (DM) is much newer subject
- We have started to scratch the surface (as a DM community)
- Many extensions are possible and necessary, to understand data phenomena in causal terms
- Interesting results on causality in DM have been obtained
- Causality in DM is related to many other DM reasoning tasks
- Maybe it is -unsurprisingly- an important underlying principle
- Or one that can lead us to a unifying concept other tasks may emerge from

