

① $y = A + Bx + Cx^2$ passing through $(1, 4)$, $(2, 7)$ and $(3, 14)$

$$4 = 1 \cdot A + 1 \cdot B + 1 \cdot C$$

$$7 = A + 2B + 4 \cdot C$$

$$14 = A + 3B + 9C$$

$$\Rightarrow A = \begin{matrix} E_1 \\ E_2 \\ E_3 \end{matrix} \begin{pmatrix} 1 & 1 & 1 & : & 4 \\ 1 & 2 & 4 & : & 7 \\ 1 & 3 & 9 & : & 14 \end{pmatrix}$$

Gaussian elimination :

$$(\bar{E}_2 - \bar{E}_1) \rightarrow (\bar{E}_2) : A = \begin{pmatrix} 1 & 1 & 1 & : & 4 \\ 0 & 1 & 3 & : & 3 \\ 1 & 3 & 9 & : & 14 \end{pmatrix}$$

$$(\bar{E}_3 - \bar{E}_1) \rightarrow (\bar{E}_3) : A = \begin{pmatrix} 1 & 1 & 1 & : & 4 \\ 0 & 1 & 3 & : & 3 \\ 0 & 2 & 8 & : & 10 \end{pmatrix}$$

$$(\bar{E}_3 - 2\bar{E}_2) \rightarrow (\bar{E}_3) : A = \begin{pmatrix} 1 & 1 & 1 & : & 4 \\ 0 & 1 & 3 & : & 3 \\ 0 & 0 & 2 & : & 4 \end{pmatrix}$$

Backward substitution:

$$C = \frac{4}{2} = 2$$

$$B = 3 - 3C = -3$$

$$A = 4 - B - C = 4 + 3 - 2 = 5$$

(2) (a)

$$A = \begin{pmatrix} 1 & 1 & -1 & : & 1 \\ 1 & 1 & 4 & : & 2 \\ 2 & -1 & 2 & : & 3 \end{pmatrix} \begin{matrix} E_1 \\ E_2 \\ E_3 \end{matrix}$$

$$(E_2 - E_1) \rightarrow (E_2)$$

$$A = \begin{pmatrix} 1 & 1 & -1 & : & 1 \\ 0 & 0 & 5 & : & 1 \\ 2 & -1 & 2 & : & 3 \end{pmatrix} \begin{matrix} E_1 \\ E_2 \\ E_3 \end{matrix} \left. \vphantom{\begin{matrix} 1 \\ 0 \\ 2 \end{matrix}} \right\} \begin{matrix} \text{need to be} \\ \text{swapped} \end{matrix}$$

(b)

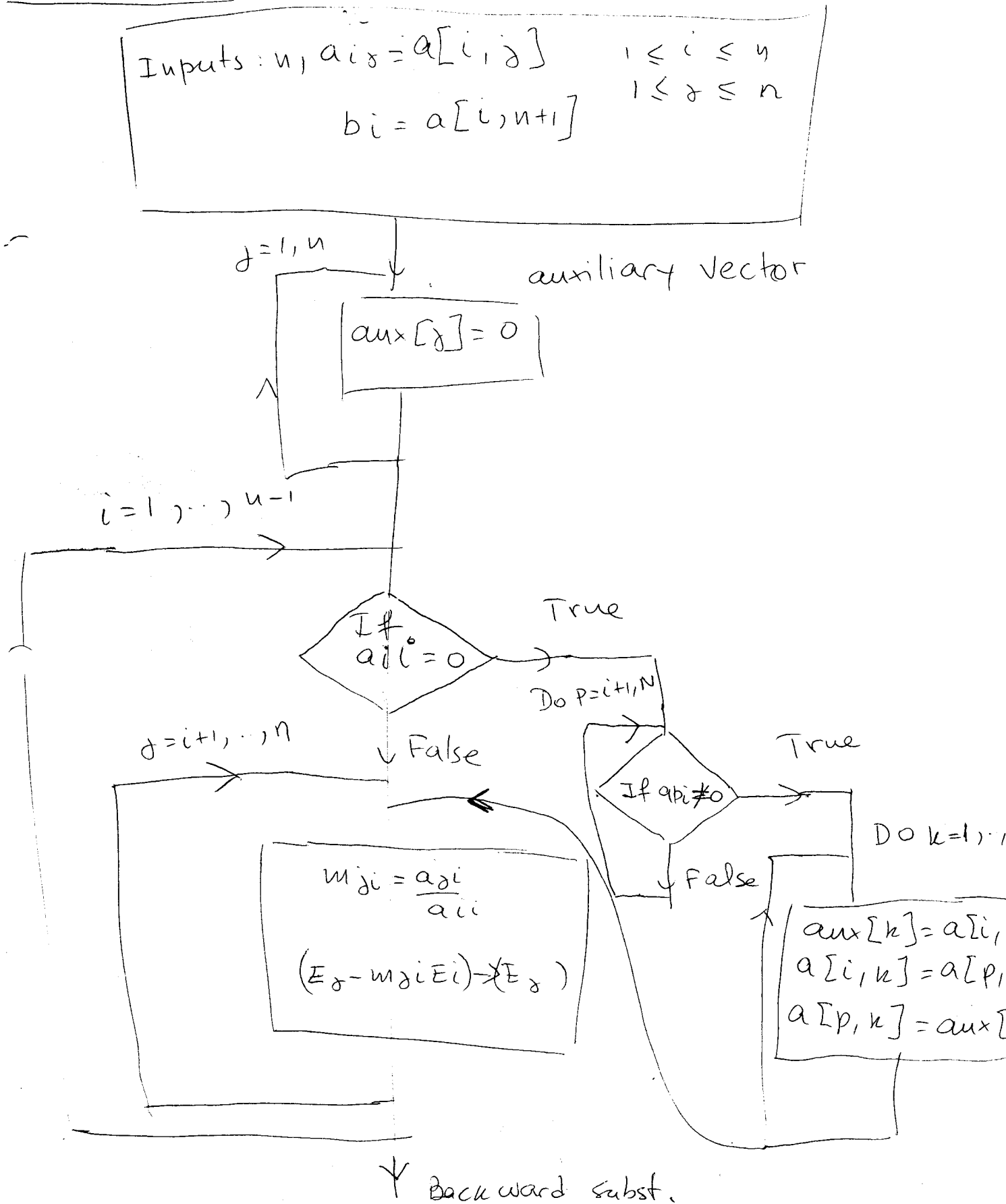
$$A = \begin{pmatrix} 0 & 1 & 1 & : & 6 \\ 1 & -2 & -3 & : & 4 \\ 1 & -1 & 1 & : & 5 \end{pmatrix} \begin{matrix} E_1 \\ E_2 \\ E_3 \end{matrix} \left. \vphantom{\begin{matrix} 0 \\ 1 \\ 1 \end{matrix}} \right\} \begin{matrix} \text{need to be} \\ \text{swapped since } a_{11} = 0 \end{matrix}$$

$$A = \begin{pmatrix} 1 & -2 & -3 & : & 4 \\ 0 & 1 & 1 & : & 6 \\ 1 & -1 & 1 & : & 5 \end{pmatrix} (E_3 - E_1) \rightarrow (E_3)$$

③ Strategy 1: Exchange rows i and p if

$$a_{ii} = 0$$

• Flow chart (Gaussian elimination)



Algorithm: Inputs: $n, a_{ij} = a[i, j] \quad 1 \leq i \leq n$
 $1 \leq j \leq n+1$

Output: x_1, \dots, x_n or message of failure

Step 1 For $j = 1, \dots, n$
 $aux[j] = 0$

Step 2 For $i = 1, \dots, n-1$ DO steps 3 - 6

Step 3 If $a_{ii} = 0$ then

Step 4 For $p = i+1, n$ (Do while $a_{pi} = 0$)

Step 5 For $k = 1, \dots, n$

$$aux[k] = a[i, k]$$

$$a[i, k] = a[p, k]$$

$$a[p, k] = aux[k]$$

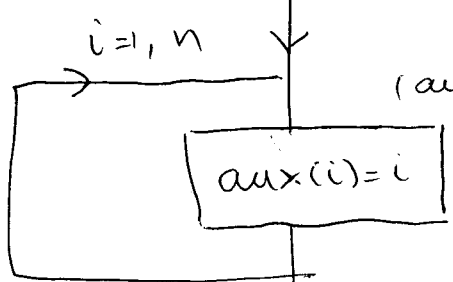
Step 6 For $j = i+1, n$

$$m_{ji} = \frac{a_{ji}}{a_{ii}}$$

$$(E_j - m_{ji} E_i) \rightarrow E_j$$

Strategy 2: Seek the largest element at the main diagonal during the Gaussian elimination

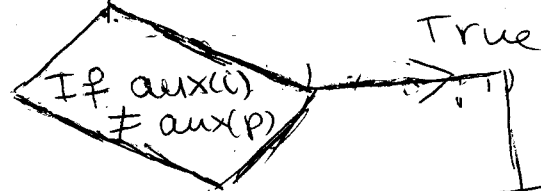
Input: M (matrix dimension); $a_{ij} = a(i, j)$, $1 \leq i \leq n$
 $1 \leq j \leq n+1$
(augmented matrix)



(auxiliary vector) this choice eliminates the need for an extra loop to interchange rows

$i = 1, n-1$

$|a(\text{aux}(p), i)| = \max_{i \leq j \leq n} |a(\text{aux}(j), i)|$
(seek maximum)



$\text{aux}(z) = \text{aux}(i)$
 $\text{aux}(i) = \text{aux}(p)$
 $\text{aux}(p) = \text{aux}(z)$

Row interchange

$j = i+1, \dots, n$

$m(\text{aux}(j), i) = a(\text{aux}(j), i) / a(\text{aux}(i), i)$
 $(E_{\text{aux}(j)} - m(\text{aux}(j), i) E_{\text{aux}(i)}) \rightarrow (E_{\text{aux}(i)})$

Backward substitution



Algorithm:

Input: number of eqs, unknowns: n
augmented matrix

Output: solution x_1, \dots, x_n or message that the linear syst. has no unique solution

Step 1: For $i=1, \dots, n$ set $\text{aux}(i) = i$

Step 2: For $i=1, \dots, n-1$ do Steps 3-6

Step 3 $|a(\text{aux}(p), i)| = \max_{i \leq j \leq n} |a(\text{aux}(j), i)|$

Step 4 If $a(\text{aux}(p), i) = 0$ then OUTPUT ('no unique solution exists')
STOP

Step 5 If $\text{aux}(i) \neq \text{aux}(p)$ then set
 $\text{aux}_2 = \text{aux}(p)$
 $\text{aux}(i) = \text{aux}(p)$
 $\text{aux}(p) = \text{aux}_2$

Step 6 For $j = i+1, \dots, n$ do Steps 7 and 8

Step 7 $m(\text{aux}(j), i) = a(\text{aux}(j), i) / a(\text{aux}(i), i)$

Step 8 $(E_{\text{aux}(j)} - m(\text{aux}(j), i) \cdot E_{\text{aux}(i)}) \rightarrow E_{\text{aux}(j)}$

Step 9 if $a(\text{aux}(n), n) = 0$ then OUTPUT ('no unique solution exists')
STOP

Backward substitution